

# Quality of Service Assurance-based Auction for Spectrum Sharing Systems

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# Introduction

## ➤ Objective:

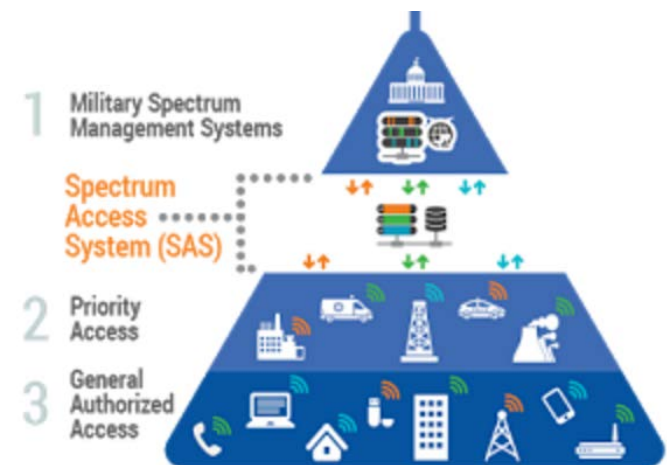
- Explore the impact of PU-SU Cooperation on spectrum opportunity management

## ➤ Proposed Solution:

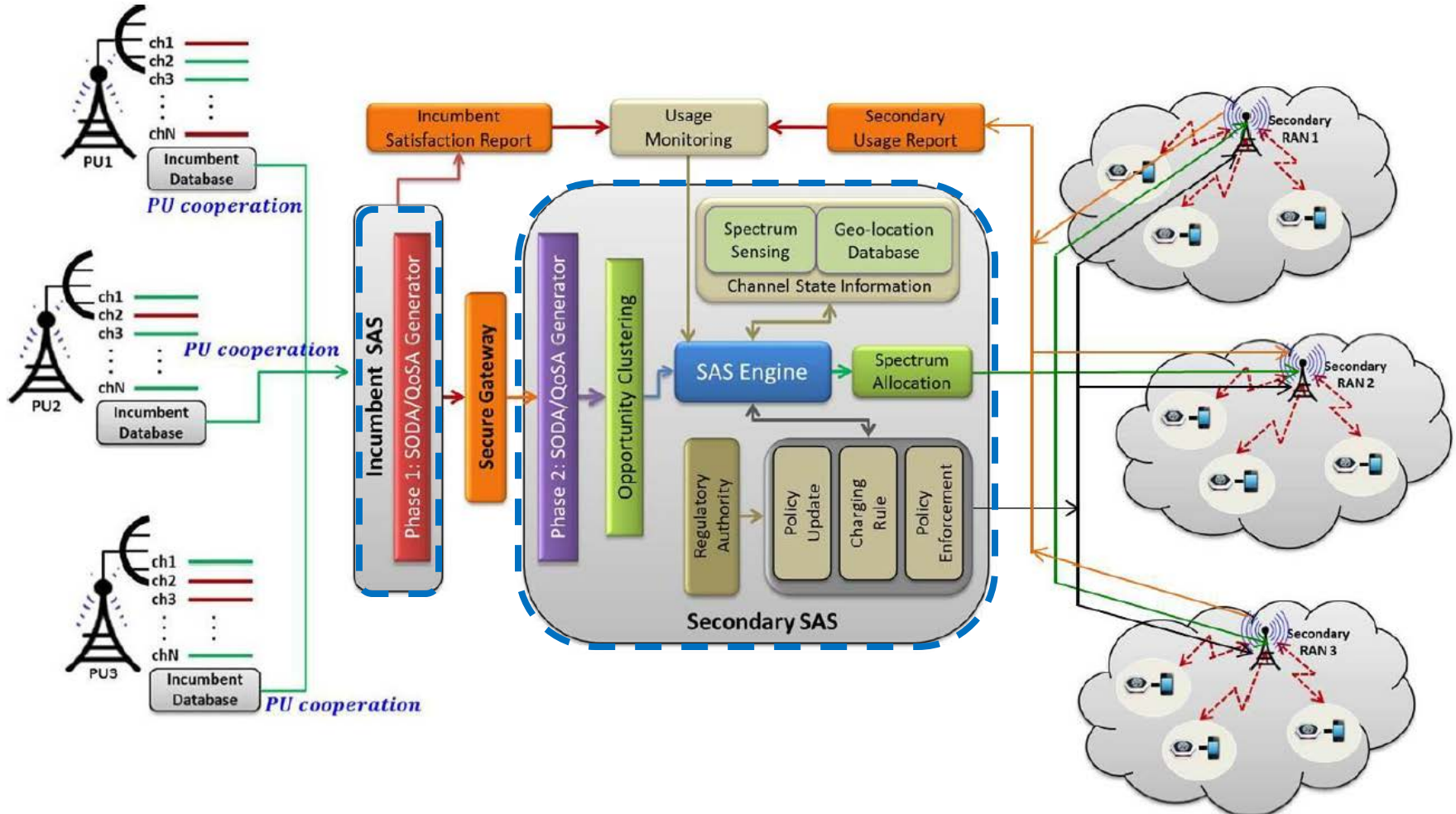
- Use QoSA values to evaluate the quality of the opportunities
  - Spectrum opportunity auctioning
  - Spectrum opportunity allocation

## ➤ Problem statement:

- Primary-Secondary Cooperation
  - Improve spectrum opportunity auctioning
  - Improve dynamic resource allocation



# PU-SU Cooperation Framework



# Opportunity Clustering with Spectrum Opportunity Duration

Duration of the opportunity

QoS Ranking		X Slots	Y Slots	...	Z Slots
	A Category	> 0.8	> 0.8		> 0.8
	B Category	0.7 – 0.8	0.7 – 0.8		0.7 – 0.8
	C Category	0.5 – 0.7	0.5 – 0.7		0.5 – 0.7
	D Category	0.3 – 0.5	0.3 – 0.5		0.3 – 0.5
	E Category	< 0.3	< 0.3		< 0.3

Effective Duration

- The normalized effective duration in terms of QoSA  $\alpha = \frac{QoS-1}{\ln(QoS)}$
- Applied to measure the quality of channel

# QoSA Calculation Procedure

## ➤ Phase-1: Incumbent SAS

- Access PU database and/or Environmental Sensing Capability (ESC) results
- Get previously stored PU usage information
- **Train the NS-HMM model parameters**

## ➤ Phase-2: Secondary SAS

- Receives estimated model parameters from I-SAS
- Uses current channel monitoring results
- **Predict the availability and duration of spectrum opportunities**



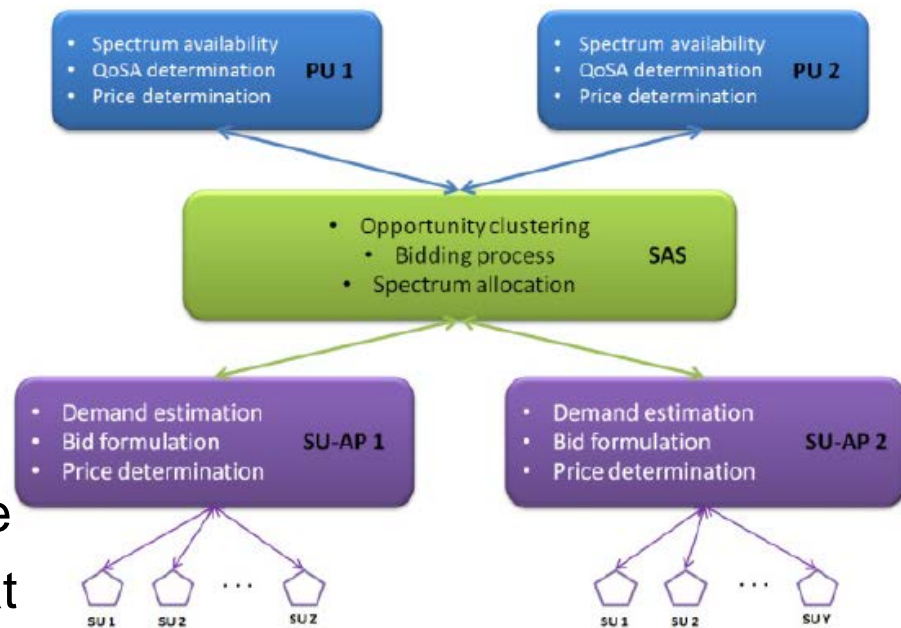
# Spectrum Opportunity Auctioning considering the Probabilistic Assurance

## ➤ An auction-based spectrum sharing framework

- Accounts for the quality of the available spectrum opportunity
  - The probabilistic assurance generated by the SAS

## ➤ PU and SU operations

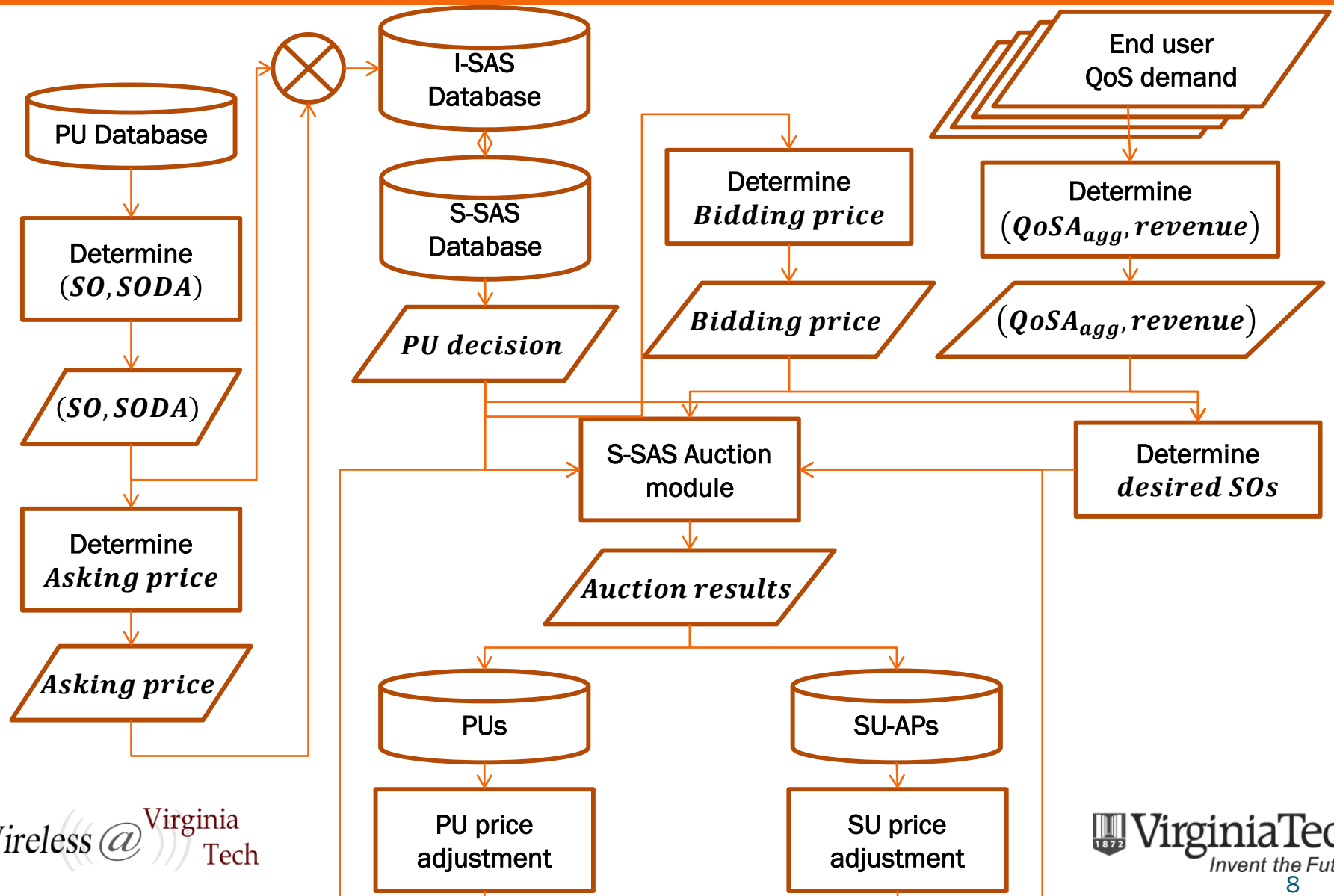
- Adjust their evaluation of the available opportunities over time
- Achieve a price combination that maximizes both the PU and SU objective



# Proposed Solution

- PUs determines the opportunities and asking prices
- I-SAS determines the SODA values for the opportunities
- SU-APs determine the aggregate QoS demand
- SU-APs determine the bidding price
- S-SAS determines the winning bids
- PUs adjust asking prices based on the auction results
- SU-APs adjusts bidding prices for unassigned channels

# Proposed Solution (Continued)





# Problem Formulation

- SU-AP tries to solve the following problem

$$\begin{aligned} \max_I \quad & P_j = \sum_{l=1}^L \sum_{i=1}^M (r_j - c_{i,j,l}) k_{i,j} \alpha_l B \tau I_{j,desired_{l,i}} \\ \text{s.t.} \quad & \sum_{l=1}^L \sum_{i=1}^M k_{i,j} \alpha_l B \tau I_{j,desired_{l,i}} \geq Q_{agg_j} \end{aligned}$$

P: SU payoffs  
L: number of opportunity clusters  
r: revenue per opportunity  
c: bidding price  
k: spectral efficiency  
 $\alpha$ : effective opportunity duration  
 $Q_{agg}$ : Aggregate QoS demand  
 $I_{desired}$ : desired opportunities

- The objective of each SU is to find the optimal bidding price per unit while
  - Maximize the payoff function (revenue per SU)
  - Satisfy the aggregated QoS demand
- Can be reduced to binary integer programming problem

# Auction Algorithm

**Goal: generate allocation indicator information for given SU bidding price**

- Spectrum opportunity clusters:  $\mathbf{SO}: \{(x_{il}\theta_{il}, \alpha_l), \forall i, l\}$ ; SU demand for each cluster:  $\mathbf{SU}_D: \{I_{jl} \in (1,0), \forall j, l\}$ ; SU bidding price for each cluster:  $\mathbf{SU}_P: \{y_{jl}, \forall j, l\}$
- Auction decision set:  $\mathbf{I}_{allocated}: \{I_{ijl} \in (1,0), \forall i, j, l\}$
- If any channel has a single claim:  
$$\mathbf{SO}: \{(x_{il}\theta_{il}, \alpha_l)\} \rightarrow \mathbf{I}_{allocated}: \{I_{ijl} \in (1,0)\}$$
- If any channel has multiple claims:  
$$\mathbf{find } j^* = \arg \max_{\forall j \in \{I_{jl}=1\}} \mathbf{SU}_P: \{y_{jl}, \forall i, l\}$$
$$\mathbf{SO}: \{(x_{il}\theta_{il}, \alpha_l)\} \rightarrow \mathbf{I}_{allocated}: \{I_{ij^*l}\}$$
$$\mathbf{0} \rightarrow \mathbf{I}_{allocated}: \{I_{ijl}\}, \forall j \in \{1,2, \dots, L; j \neq j^*\}$$
- **S-SAS forwards the allocation indicator information,**  
 $\mathbf{I}_{j,allocated}: \{I_{ijl} \in (1,0)\}$ , to each of the participating SU-APs

# Price Adjustment Algorithm

**Goal: adjust the price for given channel auction result**

- S-SAS generate list of unassigned channel and forwards  $I_{unassigned}$  to I-SAS:

*for*  $\forall l, i$ , *if*  $(SO_{il} \notin I_{allocated})$ :  $SO_{il} \rightarrow I_{unassigned}$

- I-SAS in consultation with the corresponding PU adjusts the discount factor  $\theta_{il}$

*for*  $\forall l, i \in I_{unassigned}$ ,  $\theta_{il}^{new} = \theta_{il} - \Delta_{PU_{il}}$ :  $\theta_{il}^{new} \rightarrow SO(x_{il}\theta_{il}^{new}, \alpha_l)$

- The SU-AP calculates the loss in revenue due to the current bidding price

*for*  $\forall l, i \in I_{j,desired}$ ,  $P_{jachievable} = \sum_{l=1}^L \sum_{i=1}^M (r_j - c_{i,j,l}) k_{i,j} \alpha_l B \tau I_{j,desired}$

*for*  $\forall l, i \in I_{j,allocated}$ ,  $P_{jachieved} = \sum_{l=1}^L \sum_{i=1}^M (r_j - c_{i,j,l}) k_{i,j} \alpha_l B \tau I_{j,allocated}$

$(P_{jmissd} = P_{jachievable} - P_{jachieved})$ :  $\Delta_{SU_{jl}} = \theta_{adjust} P_{jmissd}$   
 $y_{jl}^{new} = y_{jl} + \Delta_{SU_{jl}}$

- The SU-AP that gets allocated its desired channels does not perform the price adjustment mentioned in Step2

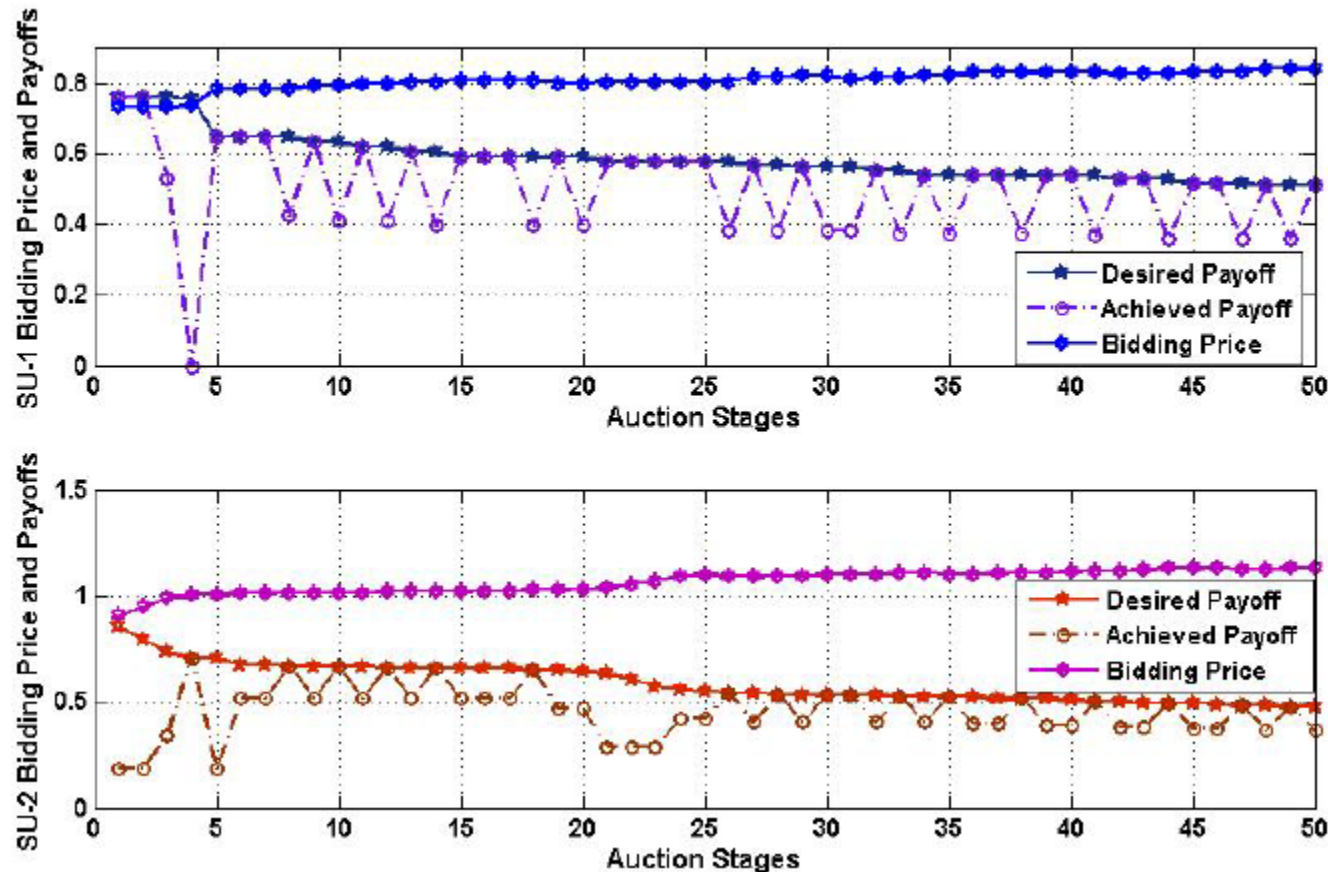
*for*  $\forall l, i$ , *if*  $(I_{desired} = I_{allocated})$ :  $\Delta_{SU_{jl}} = 0$ ;  $y_{jl}^{new} = y_{jl}$

# Simulation Parameters

Category	QoS Range	Cluster QoS	$\alpha$ value
<b>A</b>	$QoS > 0.8$	0.9	0.9491
<b>B</b>	$0.6 > QoS \geq 0.8$	0.7	0.8411
<b>C</b>	$0.4 > QoS \geq 0.6$	0.5	0.7213
<b>D</b>	$0.4 \geq QoS$	0.3	0.5814

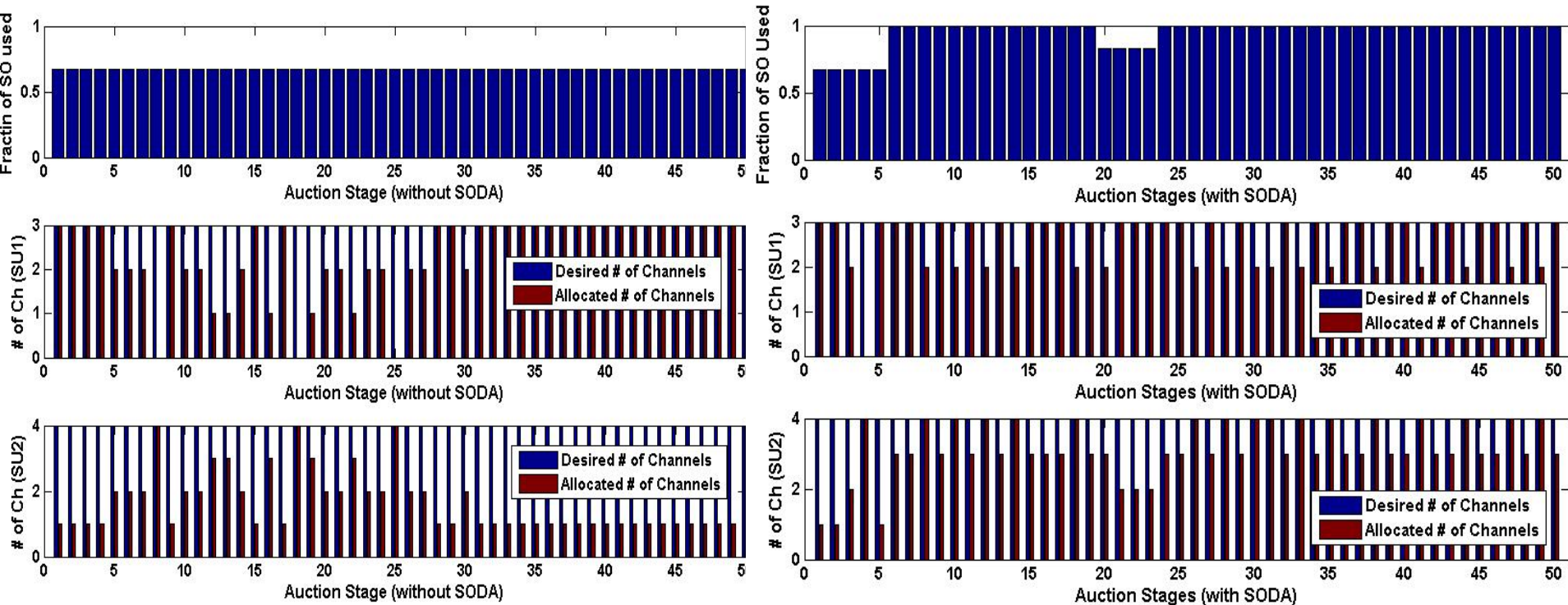
# of PUs  $M = 2$   
# of SUs  $N = 2$   
PU BW =  $10\text{MHz} \times 3$   
Duration = 1 unit  
# of cat. = 4  
Initial Price = 1 unit (PU-1)  
= 0.8 unit (PU-2)

# SU-AP Bidding Prices and Corresponding Payoffs



- The iterative adjustments in the bidding prices help the SU-APs to achieve payoffs.

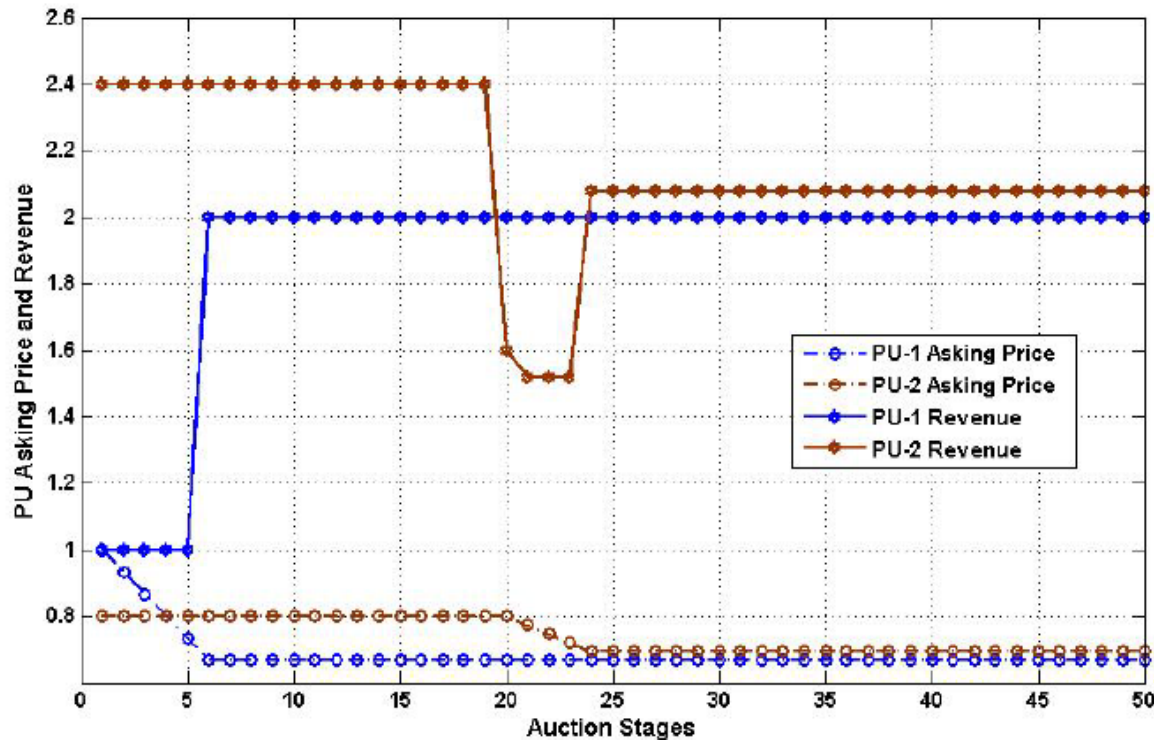
# SU-AP Channel Allocation and Overall Spectrum Usage



- The iterative price adjustment from both the PU and SU-APs helps the fair auction
- Approximately 60% of the spectrum opportunities are allocated to the SU-APs at stage 1
- All the available spectrum opportunities are allocated to the SU-Aps at later stages



# PU Asking Price and Revenues



- The revenue generated by PU1 is much lower compared to PU2 at beginning
- PU1 identifies the problem and adjusts its asking price
- The revenue earned by PU1 improves significantly by the auction stage 6
- PU2 readjusts its asking price at stage 20

# Conclusion and Future Works

- We proposed an auction-based spectrum sharing framework includes the spectrum opportunity duration information
- We proposed an iterative auction algorithm that quality of the spectrum opportunities is decided
- The proposed approach improved the percentile of the available spectrum opportunities used by the SUs
- SODA values allow estimation of effective opportunity duration
  - Achievable QoS from the opportunities can be determined
- SUs request opportunities to meet their QoS demand
  - Translating QoS demand for dedicated channel to equivalent group of opportunities